

# Specifying Teletype Behavior for Automated Handling of Exercises on Interaktive Haskell Programs

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# Motivation

- ▶ We teach Haskell in a course on programming paradigms
- ▶ Each week students solve small programming tasks
- ▶ We want to automatically test submissions for correctness
  - ▶ Easy for pure functions (QuickCheck etc.)
  - ▶ But there is also IO
- ▶ Goal:
  - ▶ Flexible and easy to use way to specify and automatically test Haskell IO programs
  - ▶ Enable further optimization and analysis

# Teletype IO Programs

- ▶ Two basic operations:

```
readLn :: Read a => IO a
print  :: Show a => a -> IO ()
```

- ▶ Results of IO computations do not contain information on side effects
- ▶ For programs with type `IO ()`, comparing against the result only checks termination
- ▶ Therefore we capture a programs monadic effects with traces over basic operations

# Current Solution - Algebraic Effects

## Swiestra and Altenkirch, Haskell Workshop 2007:

“By constructing an internal model of such concepts [teletype IO, ...] within our programming language, we can test, debug and reason about programs that perform IO as if they were pure.”

```
data Teletype a = ReadLine (String -> Teletype a)
                | WriteLine String (Teletype a)
                | Return a

instance Monad Teletype

readLn :: Read a => Teletype a
print  :: Show a => a -> Teletype ()

runTeletype :: Teletype a -> [String] -> Trace a

data Trace a = Read String (Trace a)
              | Write String (Trace a)
              | Result a
```

# Current Solution - Implementing Tasks

(Silently) hand students this verison of IO

```
type IO = Teletype

type Input = String {- or somthing else -}

validInputs :: Gen [String]
checkCorrectness :: Trace -> Bool
correctTrace :: [Input] -> Trace {- maybe from sample solution -}

prog :: IO () {- student solution -}
check = quickCheck $ forAll validInputs
  \xs -> counterexample (show $ correctTrace xs) $
  checkCorrectness (runTeletype prog xs)
```

For each task implement `validInputs`, `checkCorrectness` and `correctTrace`

# Current Solution - Problems

- ▶ Implementing new tasks is hard
  - ▶ especially checking correctness
  - ▶ and accounting for optionality
- ▶ Inconsistencies between components possible
- ▶ explicit (ad hoc) checking instead of declarative description
- ▶ Automatization beyond testing not possible (e.g. task generation)

# Specification Language

- ▶ Build sufficiently expressive specification language
  - ▶ Describe basic IO behavior
  - ▶ Leave room for optional extra information (e.g. additional output)
- ▶ Automatically derive components for testing
- ▶ Keep design minimal to enable more automatization and analysis in the future
  - ▶ Task generation
  - ▶ Deriving (idiomatic) sample solutions
  - ▶ Generate helpful feedback

## Example specification

*Read a natural number  $n$  from `stdin`, then read  $n$  additional numbers and print the sum of those  $n$  numbers to `stdout`.*



# Example specification

*Read* a natural number  $n$  from *stdin*, then *read*  $n$  additional numbers and *print the sum of those  $n$  numbers* to *stdout*.

- ▶ primitive operations
- ▶ flow control & iteration
- ▶ backwards references/variables

# Primitive Operations

- ▶ Read a value of “type”  $\tau$  into  $x$ :  $[\triangleright x]^\tau$
- ▶ Print the value of a term  $t$ :  $[t \triangleright]$ 
  - ▶ Term language can vary depending on what kind of expressions are needed

# Flow Control & Iteration

- ▶ Branch based on a predicate  $p$ :  $s_1 \triangleleft p \triangleright s_2$
- ▶ Iterate (sub-)specification:  $s \rightarrow^E$
- ▶ Terminate iteration: **E**

# Variables

- ▶ With conventional variables, iterating over reads shadows references to old values
- ▶ Therefore **all** variables store **lists!**
  - ▶ Read = append new value to list
  - ▶ No value is ever overwritten
  - ▶ Two access methods for each variable
    - ▶ Last/current value:  $x_C$
    - ▶ All values:  $x_A$

# Example Task

*Read a natural number  $n$  from `stdin`, then read  $n$  additional numbers and print the sum of those  $n$  numbers to `stdout`.*

$[\triangleright n]^{\mathbb{N}} \cdot ([\triangleright x]^{\mathbb{Z}} \angle \text{len}(x_A) = n_C \triangleright \mathbf{E})^{\rightarrow \mathbf{E}} \cdot [\{\text{sum}(x_A)\} \triangleright]$

# Optionality

Read a natural number  $n$  from *stdin*, then read  $n$  additional numbers and print the sum of those  $n$  numbers to *stdout*. **The program might print each time how many more summands it is still expecting.**

Extend writes to specify a set of possible outputs

$$[\{t_1, t_2, \dots\} \triangleright]$$

Encode optionality / no output with special “empty” term  $\varepsilon$

$$[\triangleright n]^{\mathbb{N}}([\{\varepsilon, n_C - \text{len}(x_A)\} \triangleright][\triangleright x]^{\mathbb{Z}} \angle \text{len}(x_A) = n_C \triangleright \mathbf{E}) \rightarrow^{\mathbf{E}}[\{\text{sum}(x_A)\} \triangleright]$$

# Testing Behavior

- ▶ Given program  $p$  and specification  $s$
- ▶ Does  $p$  match the behavior described by  $s$ ?
  1. Define what it means for a single program run/trace to match the specification
  2. Does this hold for all traces  $p$  can produce?
  3. If yes, then  $p$  fulfills the specification

# Example

- ▶ Trace: ?2 ?5 ?3 !8
- ▶ Match against specification by left-to-right traversal

$[\triangleright n]^{\mathbb{N}}$	$([\{\varepsilon, n_C - \text{len}(x_A)\} \triangleright])$	$[\triangleright x]^{\mathbb{Z}}$	$\angle \text{len}(x_A) = n_C \triangleright \mathbf{E} \rightarrow^{\mathbf{E}}$	$[\{\text{sum}(x_A)\} \triangleright]$
$2 \rightarrow n$	$\varepsilon$	$5 \rightarrow n$	$\text{len}(\text{Nil}) = 2 : \text{False}$	
	$\varepsilon$	$3 \rightarrow n$	$\text{len}([5]) = 2 : \text{False}$	
			$\text{len}([5, 3]) = 2 : \text{True}$	$8 \in \{\text{sum}([5, 3])\}$

- ▶ Successful match since the trace is consumed completely



# Automatic Testing

- ▶ Remember that we need at least the following components

```
validInputs :: Gen [Input]
checkCorrectness :: Trace -> Bool
```

- ▶ We can reverse trace matching to sample valid program runs from specifications
- ▶ Such a run contains a valid `Input` sequence
- ▶ Generalizing runs w.r.t. optionality lets us check if a run of an actual program is contained in the general case
- ▶ If this is not the case we found a counterexample
- ▶ So the program does not have the described behavior

# Example

$$[\triangleright n]^{\mathbb{N}}([\{\varepsilon, n_C - \text{len}(x_A)\} \triangleright][\triangleright x]^{\mathbb{Z}} \angle \text{len}(x_A) = n_C \triangleright \mathbf{E}) \xrightarrow{\mathbf{E}} [\{\text{sum}(x_A)\} \triangleright]$$

Sample trace:  $\{3\} \{\varepsilon, 3\} \{1\} \{\varepsilon, 2\} \{7\} \{\varepsilon, 1\} \{4\} \{12\}$

Matching:

▶  $\{3\} \{3\} \{1\} \{2\} \{7\} \{1\} \{4\} \{12\}$

▶  $\{3\} \{3\} \{1\} \{7\} \{1\} \{4\} \{12\}$

Non-Matching:  $\{3\} \{4\} \{1\} \{3\} \{7\} \{2\} \{4\} \{12\}$

# EDSL Implementation

## Embedding specifications in Haskell

```
exampleTask :: Specification
exampleTask =
  readInput "n" NatTy <>
  tillE (
    branch
      ((\xs n -> length xs == n) <$> getAll @Int "x" <*> getCurrent "n")
      (readInput "x" IntTy)
      e
  ) <>
  writeOutput [sum <$> getAll @Int "x"]

test :: IOtt () -> IO ()
test prog = quickCheck $ prog `fulfills` exampleTask
```

$$[\triangleright n]^{\mathbb{N}} \cdot ([\triangleright x]^{\mathbb{Z}} \angle len(x_A) = n_C \triangleright \mathbf{E})^{\rightarrow^{\mathbf{E}}} \cdot [\{sum(x_A)\} \triangleright]$$

# Conclusion

- ▶ Simple framework for testing teletype IO behavior
- ▶ Clear and declarative specifications
- ▶ Potential uses beyond testing:
  - ▶ Provide better feedback
  - ▶ Specifications can be randomly generated
  - ▶ Automatic generation of (idiomatic) sample solutions
  - ▶ Maybe even automatic verbalization (of tasks)