$\begin{array}{c} \text{SS 2019} \\ \text{Sheet 1} \\ 30^{\text{th}} \text{ April 2019} \end{array}$ 

# Exercise: "Modelling of concurrent systems"

#### Submission instructions:

- (a) Please submit your solutions (containing the sheet number and matriculate number) by 26.04.19 by 17:00 hours.
- (b) You can either hand it over in person after the lecture on 26<sup>th</sup> April or send it to me by email: harsh.beohar@uni-due.de.
- (c) The tutorial will take place on 30.04.19 from 14:15–15:30 in Room LF 035.

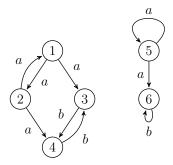
#### Task 1 Mutual exclusion

Consider two processes  $P_1$  and  $P_2$  that share a common resource R. The behaviour of the two processes is identical; both can be in one of the following states: neutral (doing nothing), trying to access the resource R, or critical section (i.e., accessing the resource R). Both the processes are initially set to neutral state. Once the processes are in neutral states, they can try to enter the critical section. However, there is the following restriction on how to enter the critical section: if the process  $P_i$  (for  $i \in \{1, 2\}$ ) is in trying state and  $P_j$  (for  $j \in \{1, 2\}$  and  $j \neq i$ ) is not in the critical section, then only  $P_i$  can enter the critical section. Thus, our task is to ensure that both the processes do not simultaneously access the resource R.

- (a) Model the above situation using a labelled transition system. Please choose the labelling of states and actions in a sensible way that is relevant to the above context. Also, give an explanation to each of the labels used by you in your model. (Points: 3)
- (b) Based on your above model, answer the following queries with explanations.
  - i) Does your model guarantees the mutual exclusive access of the resource *R*? (Points: 1)
  - ii) Are there any unfair situations in your model, where one process does not allow the other process to enter the critical section forever? How many such unfair situations are present in your model? (Points: 2)

#### Task 2 Language Equivalence

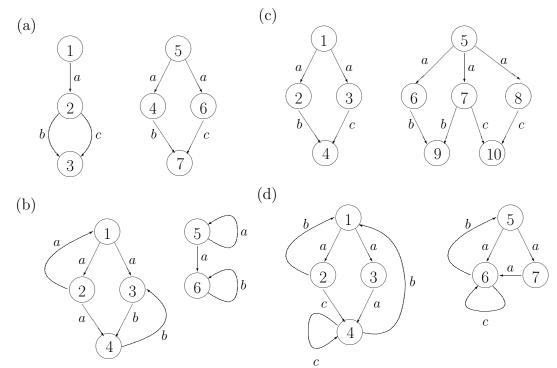
Consider the following labelled transition system:



- (a) Using the *regular expressions*, give the set of all finite words generated by the states 1, 5. (Points: 2)
- (b) Give the set of all infinite words generated by the states 1, 5. (Points: 2)
- (c) Does the states 1,5 generate the same set of (in)finite words? (Points: 2)

### <u>Task 3</u> Behavioural equivalences

Consider the following pair of labelled transition systems with the alphabet  $L = \{a, b, c\}$ . In each of the pair, establish whether the states 1 and 5 are failure equivalent and bisimilar? Explain your answers and compute the relation explicitly in case of bisimulation using the partition refinement algorithm. (**Points: 8**)



<u>Task 4</u> Simulation equivalence

Let  $(S, L, \rightarrow)$  be an labelled transition system. A binary relation  $R \subseteq S \times S$  is a *simulation* relation if and only if the following condition is satisfied.

$$\forall_{s,t,s'\in S} \forall_{a\in L} \ \Big( (sRt \wedge s \xrightarrow{a} s') \implies \exists_{t'\in S} \ (t \xrightarrow{a} t' \wedge s'Rt') \Big).$$

A state t simulates the behaviour of another state s, denoted  $s \leq t$ , if there exists a simulation relation R such that sRt.

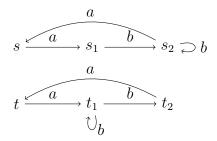
- (a) Show that the relation  $\leq$  is a preorder on the set of states S. (Points: 2)
- (b) Consider the *simulation equivalence* on the set of states S defined in the following way:  $s \approx t \iff s \leq t \land t \leq s$ . In other words, two states s and t are simulation equivalent if and only if t simulates the behaviour of s and vice versa.

Does simulation equivalence coincides with bisimilarity, i.e., whether is it the case that for any two states s, t we have s and t are bisimilar if and only if s and t are simulation equivalent? If yes, then prove it. If not, give an example supporting your claim. (Points: 2)

## <u>Task 5</u> Comparitive concurrency semantics

Prove the following facts (Points: 2+2+2):

- (a) If two states of a transition system are bisimilar then they are failure equivalent.
- (b) Show that finitary trace equivalence coincides with bisimulation under the assumption of determinism.
- (c) Consider the following transition systems from the state s, t as depicted below: Are the two states s and t bisimilar? If yes, argue why they satisfy the same set



of HML formulas. Otherwise give a formula expressible in HML that distinguishes the two upto bisimilarity. (Points: 1+1)